

# Infinite Memory Engine (IME)

## Partial Non-Deterministic I/O System for Exascale

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# Who's IME?

**Lead Architect: Paul Nowoczynski**

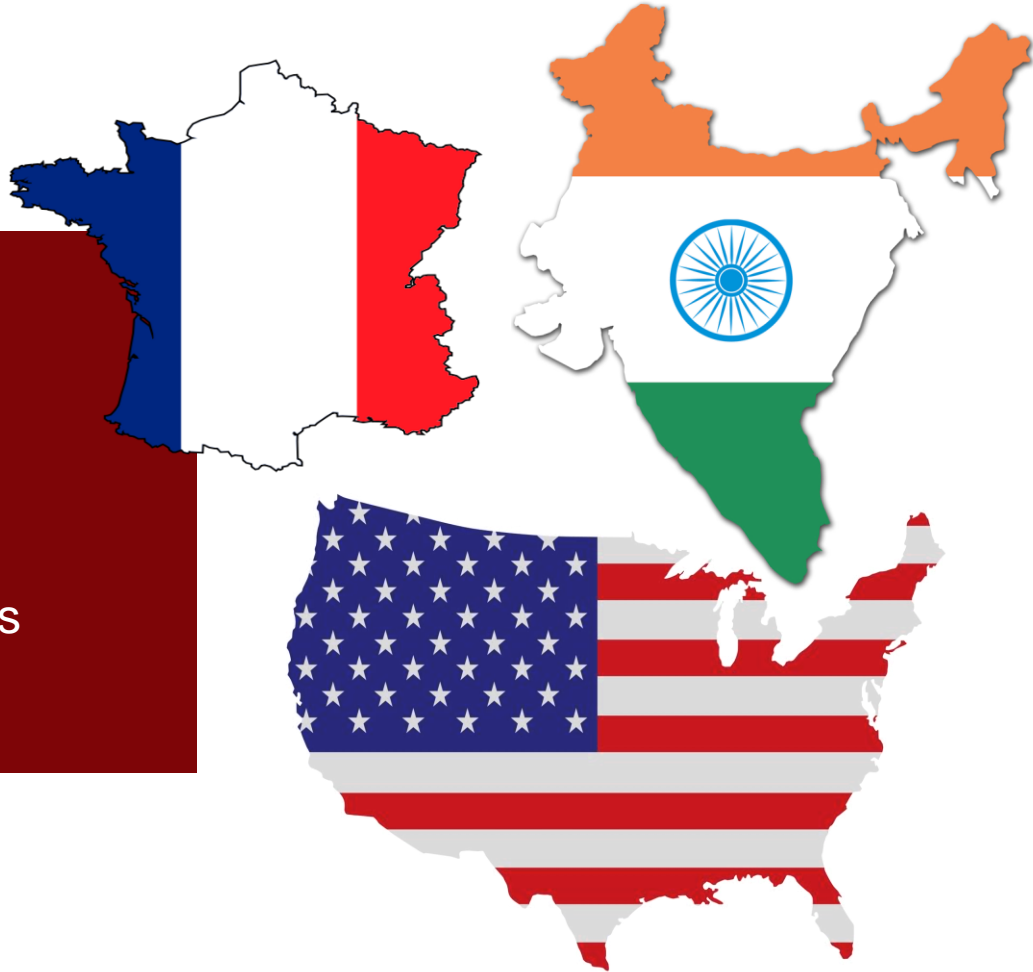
**R&D team on 3 continents**

▮ Maryland, US

▮ Paris, France

▮ Pune, India

First project developed in the new Paris  
Advanced Technical Center



# Why IME?

## Bandwidth coarse estimation

**Bandwidth needs next-gen pre-Exascale systems**

**Rules of thumb**

**1/ Checkpointing less than 6 minutes per hours**

**2/ Checkpointing means draining half of system memory**

**Pre-Exascale system:**

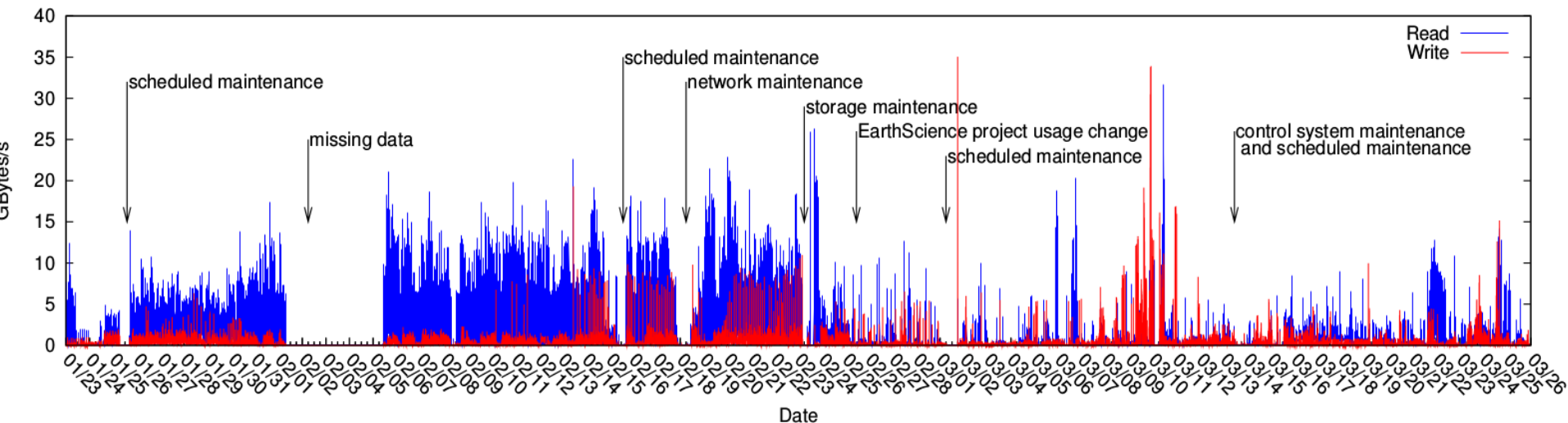
**4 Petabyte → bandwidth requirement 5.6 TB/s**

**Oakridge,**

**Teng Wang, Weikuan Yu et al. “ An Efficient Distributed Burst Buffer for Linux”**

# Why IME?

## Bandwidth detailed view



**99% of the time the IO sub-system is stressed below 30% of its bandwidth**  
**70% of the time the system is stress under 5% of its peak bandwidth**

Argonne lab.

P. Carns, K. Harms et al. *Understanding and Improving Computational Science Storage Access through Continuous Characterization*

# Why IME? economics

SSD reshuffle the parameters

Latency / 40 : 4ms  $\rightarrow$  0,1 ms

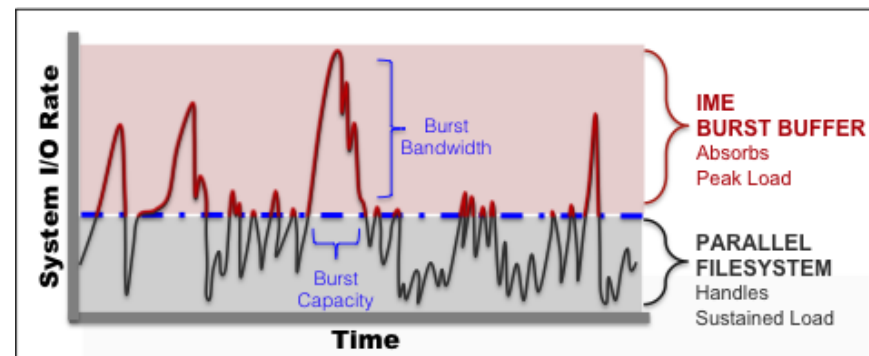
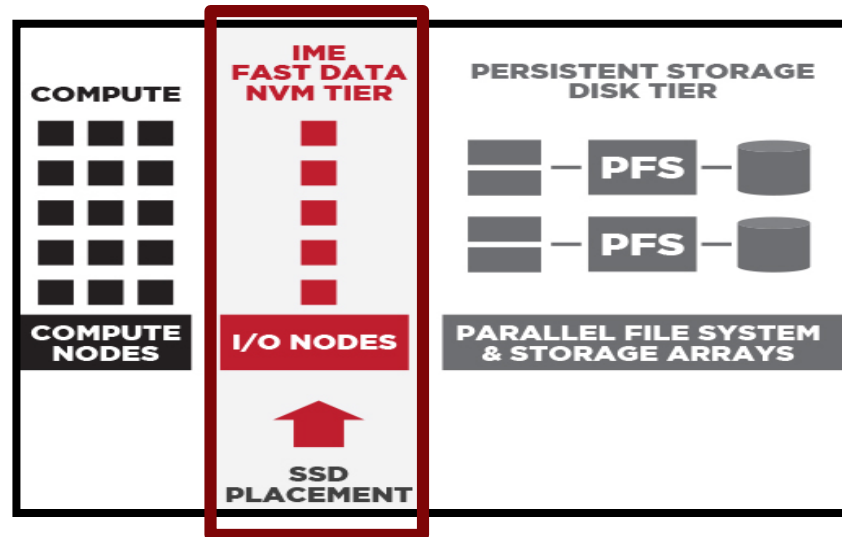
Bandwidth x 3: 150  $\rightarrow$  450 MB/s

Capacity / 8 : 8  $\rightarrow$  1TB.

Cost x 10 \$ 0,05/Gbit  $\rightarrow$  \$0.04

+ SSD deprecation rate !!

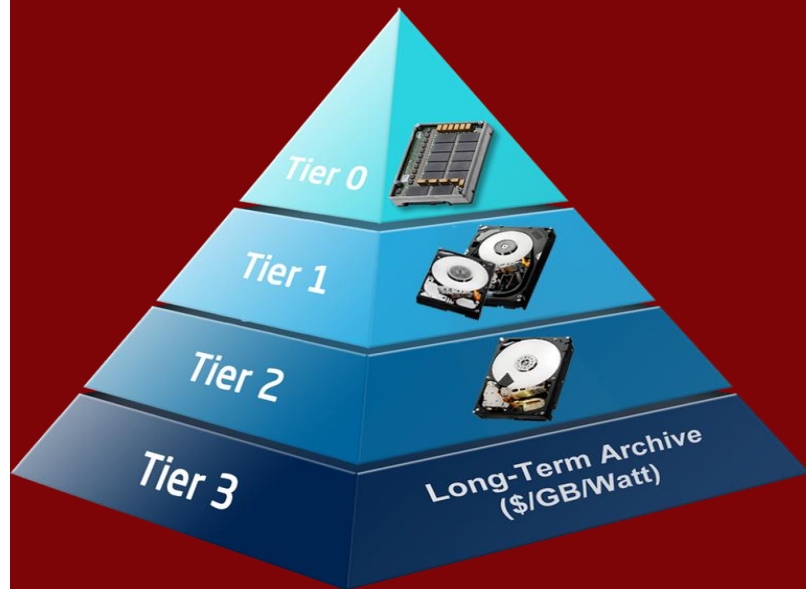
What can we do with a costly high  
bandwidth low latency technology ?



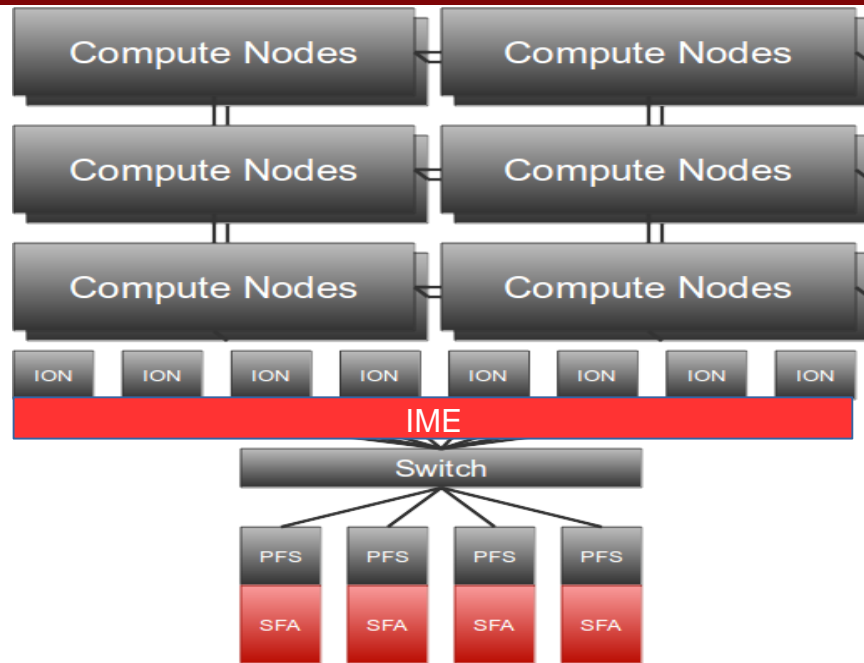
# What is IME?

Distributed virtually shared coherent array of SSDs

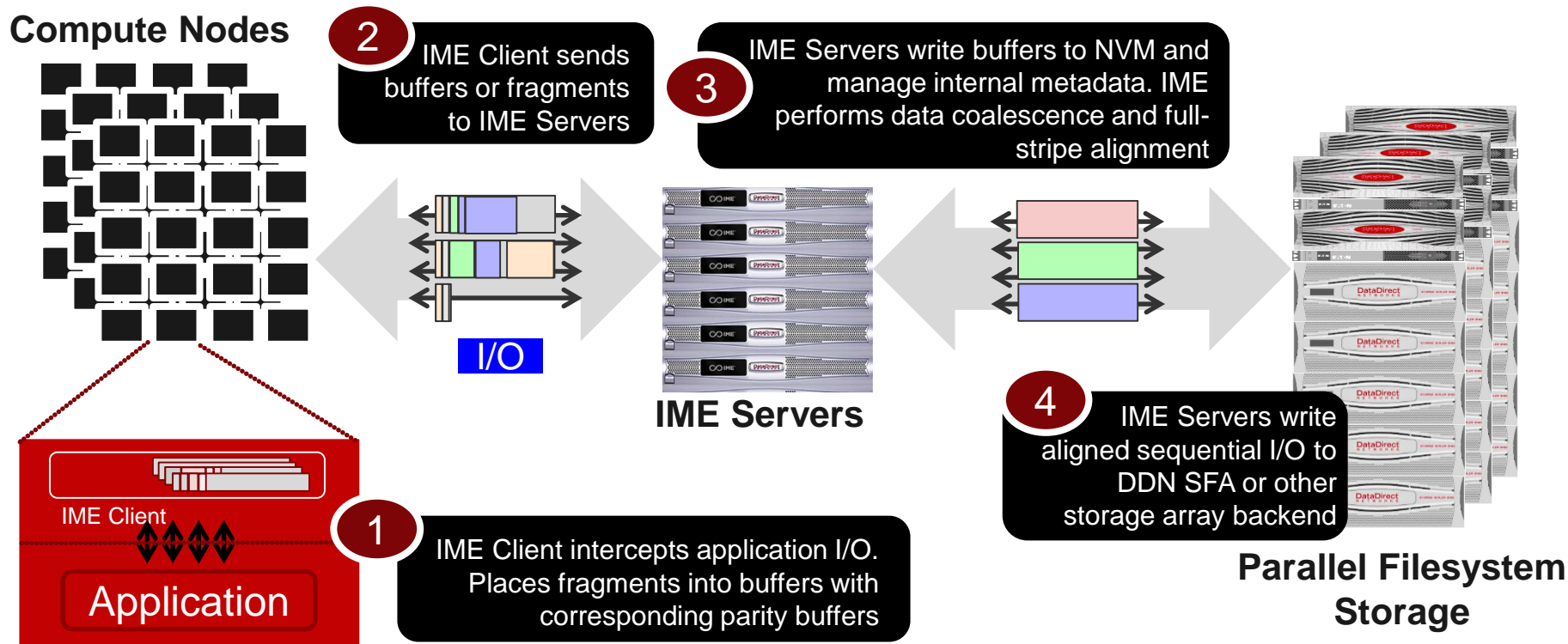
## Hardware tiering



## System architecture



# IME New Data Flow, New software stack



# IME means Scalability and Reliability

Alleviate lock pressure: byte addressable, no page lock

**Every IO is recorded as a frag:**

**File ID**

**Address**

**Length**

**Frag are aggregated in large network buffer**

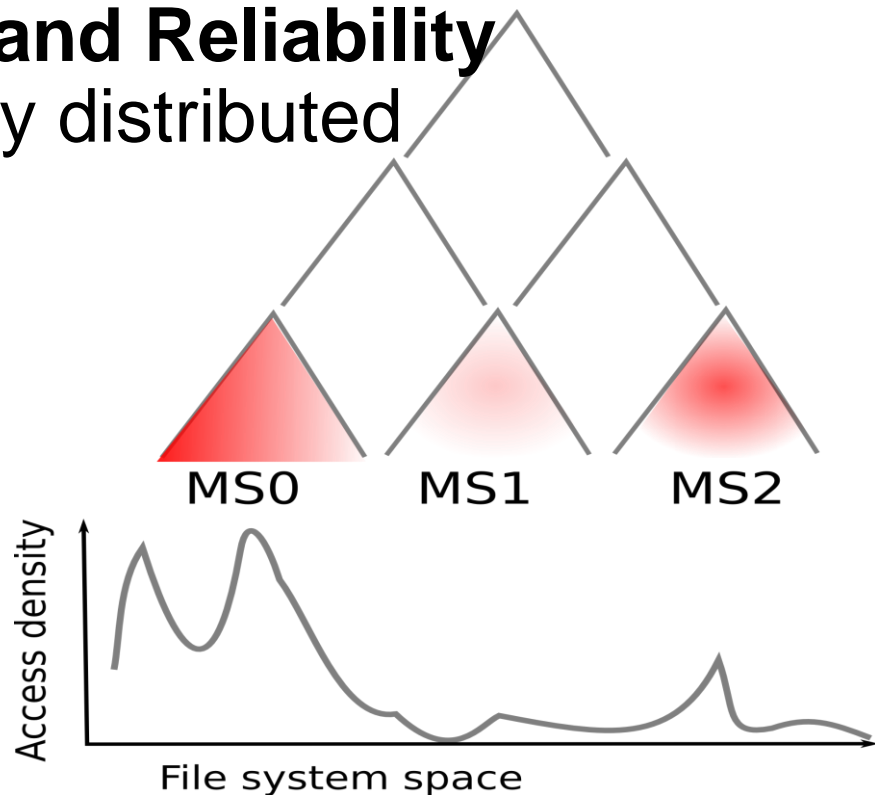
**Frag are weakly ordered, weak Posix compliance**



# IME means Scalability and Reliability

## Metadata directory is fully distributed

**No centralized directory**  
**Distribution independent from structure**  
→ no hot directory issue  
**Common knowledge ensured through hash**  
→ CRUSH like hash function

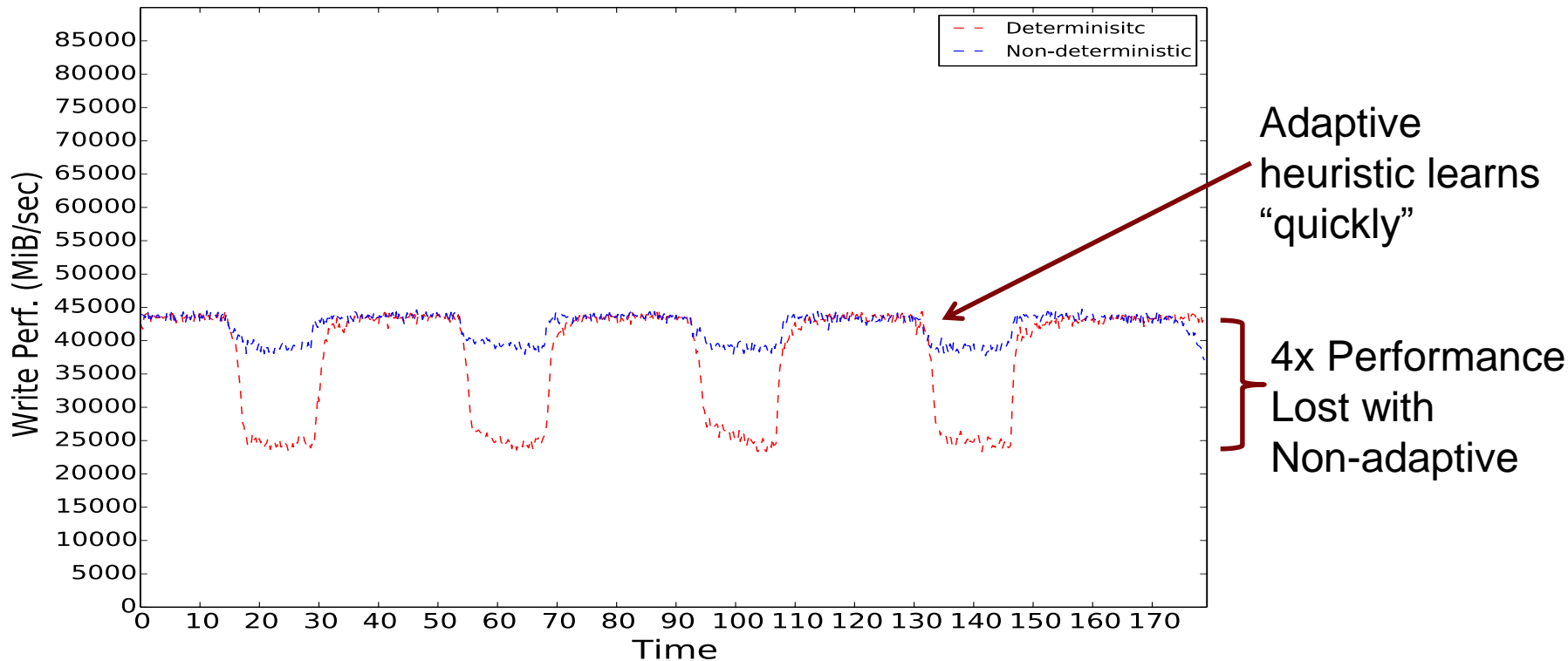


Santa Cruz U., Supercomputing 2006

Sage Weil, Scott Brandt et al. *CRUSH: Controlled, Scalable, Decentralized Placement of Replicated Data*

# From fault tolerance to resilience

## Leverage server fault tolerance for load balancing



# IME means Scalability and Reliability

Parity and data protection is not going to scale

**IME software only**

**Erasur coding not scalable on server side**

**Done on the IME client side**

- **Vector instructions**
- **Heavily multithreaded**

Pittsburgh Supercomputing Center., Supercomputing 2006  
Paul Nowoczynski et al. Zest Checkpoint Storage System for Large Supercomputers

# IME Fault Recovery

- SSD Failure

- Data recovery responsibility belongs to IME server with failed device
- Recovered data are rewritten to remaining SSDs

- IME server failure

- Recovery performed by remaining servers
  - Aka “Node Ejection”
- Fully declustered
  - Remaining servers share rebuild processing

# IME Boost

Accelerates applications

→ especially small or mal-aligned I/O

No page lock

Client coalescing:

→ Bulk RPC transfer

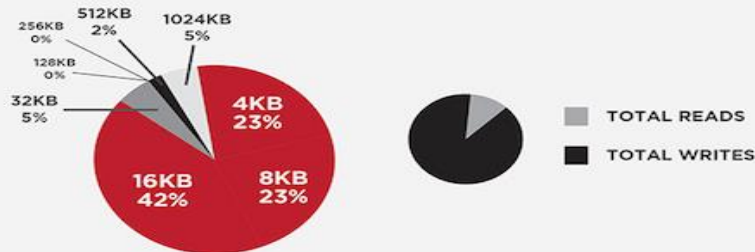
→ Save SSD write cycle

Server flush scheme

→ tune for PFS parameters

## WRITE DISTRIBUTION FOR MULTI-DISCIPLINARY HPC CLUSTER

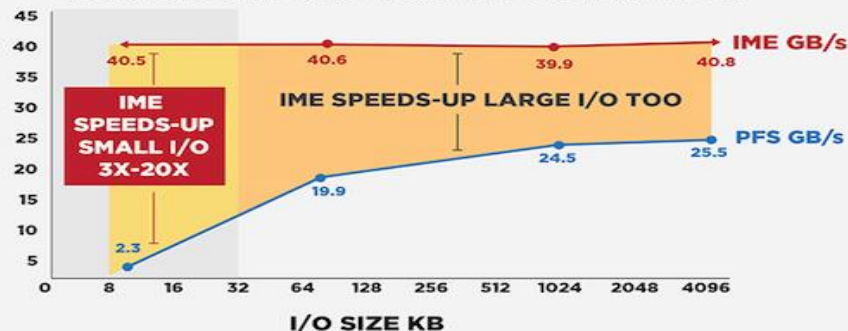
EVEN LARGE HPC SITES DRIVE A LOT OF SMALL I/O



90% OF ALL I/O IN TYPICAL HPC DATACENTERS IS <32KB IN SIZE

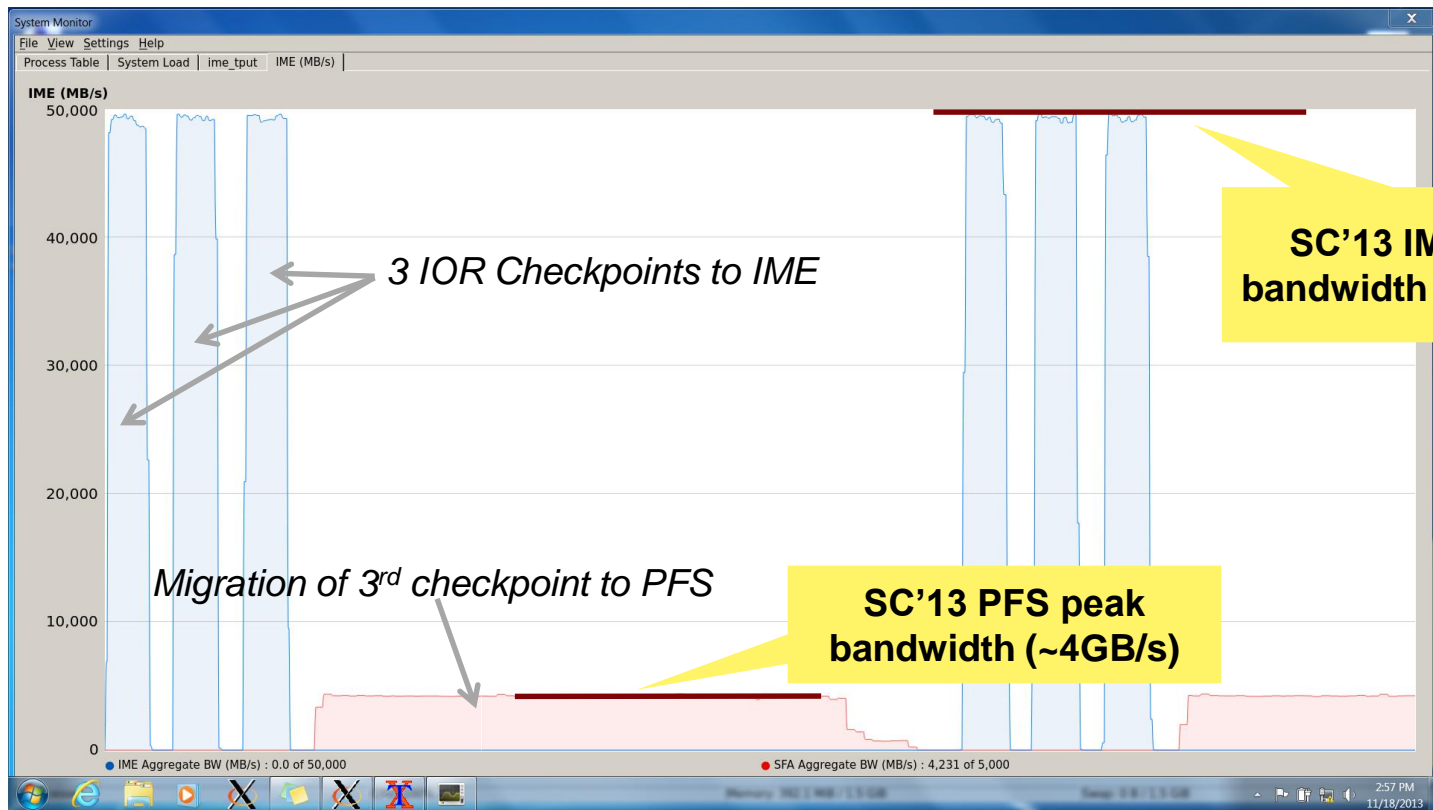
## HOW IME HELPS

ENABLES HIGHER PEAK BANDWIDTH THAN  
DISK-BASED PFS, ESPECIALLY FOR SMALL I/Os



# IME Boost

## IOR SSF 4k interleaved writes (SC'13)

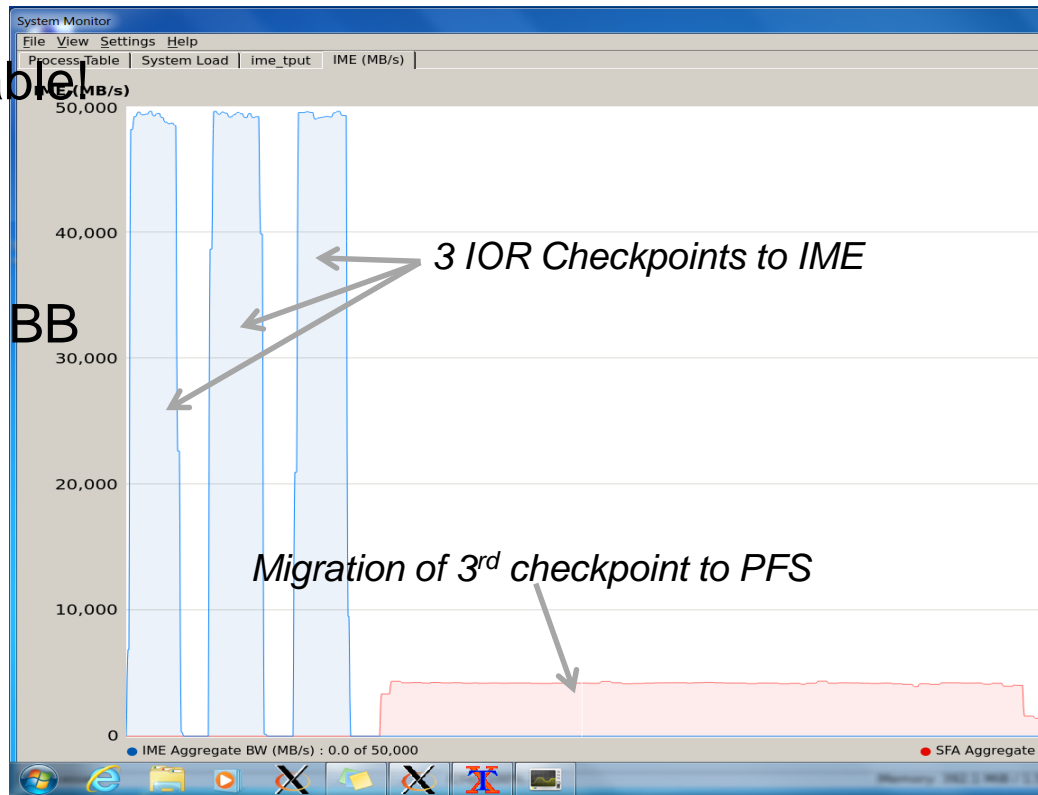


# IME Boost

## Reduce data movement

Copying all data to PFS is not desirable

- Reduce data movement
- 3<sup>rd</sup> checkpoint should be freed from BB following copy to PFS
- ***Explicit management is required***



# Future: Node Local NVM

Node local provides a new set of challenges.

- Affiliation of file components to specific compute nodes must be expressed
- Will a job actually run on the compute nodes where data has been staged-in?
  - At large issue in HPC deployments
- Blur the distinction between client and server.
  - Performance instructiveness in compute node



# IME take away

- Extra tier of SSD in storage hierarchy
  - Re-design the software stack to address real issues
  - Resilience and scalability
- Keep data as close as possible from compute node
- Pluggable with other IO storage technology
  - IO libraries
  - Object storage
- Pave the way for future evolutions
  - A component of the Exascale storage stack



# Thanks !

# IME Approach

## Distributed Hash Table + Log Structuring

